

Wilf-equivalence of partial permutations

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Joint work in part with [Tian Han](#), [Sergey Kitaev](#), and [Philip Zhang](#)



1 Definitions

2 Results

3 Conjectures

Partial permutations

Hamaker (PP'24): A **partial permutation** of size m is a pair of m -tuples of distinct positive integers (A, B) , that is

$$A = (a_1, \dots, a_m), \quad B = (b_1, \dots, b_m).$$

A k -extension of a partial permutation (A, B) of size j is a permutation $\sigma \in S_k$ such that

$$\sigma(a_j) = b_j \quad \text{for each } j = 1, \dots, m.$$

This is defined for $k \geq m$.

A k -completion $(A, B)_k$ of a partial permutation (A, B) is the set of all k -extensions of (A, B) .

For example,

- $((2, 4), (1, 2))_4 = \{3142, 4132\}$,
- $((2), (1))_4 = \{2134, 2143, 3124, 3142, 4123, 4132\}$,
- $(A, B)_k = \emptyset$ for $k < \max(A \cup B)$,
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Pattern avoidance

- A permutation π contains an occurrence (or instance) of pattern σ , there is a subsequence of π order-isomorphic to σ .
- π avoids σ if π does not contain an occurrence of σ .
- π avoids a set of patterns S if π avoids every pattern in $\sigma \in S$.
- Denote the set of permutations of size n avoiding a pattern σ (resp. a set of patterns S) by $\text{Av}_n(\sigma)$ (resp. by $\text{Av}_n(S)$).
- Call sets of patterns S and T Wilf-equivalent if $|\text{Av}_n(S)| = |\text{Av}_n(T)|$ for all $n \geq 0$, and denote this by $S \sim T$.

Wilf-equivalence for partial permutations

- Call partial permutations (A, B) and (C, D) **k -Wilf-equivalent** if $(A, B)_k \sim (C, D)_k$.
- Call partial permutations (A, B) and (C, D) Wilf-equivalent if $(A, B)_k \sim (C, D)_k$ for all $k \geq \max(A \cup B \cup C \cup D)$. Notation: $(A, B) \sim (C, D)$.
- For example, the following partial permutations are 3-Wilf-equivalent:

$$(13, 13)_3 = (12, 12)_3 \sim (13, 12)_3 \sim (23, 12)_3, \quad \text{i.e. } 123 \sim 132 \sim 312,$$

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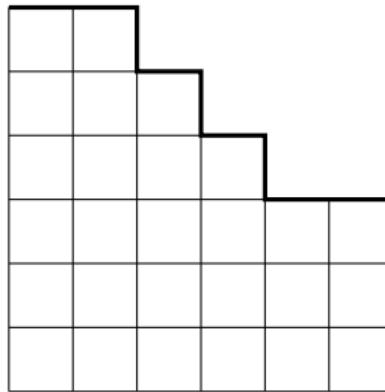
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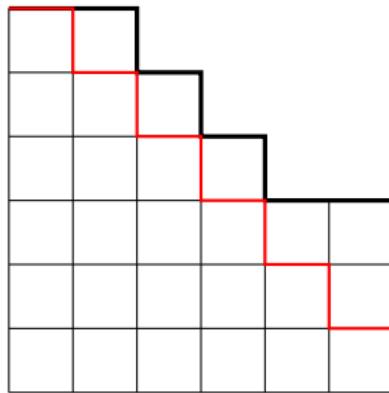
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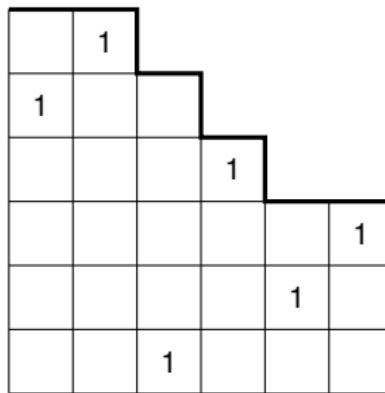
(NE)-shape-Wilf equivalence



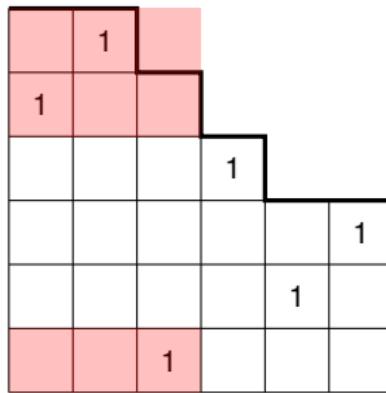
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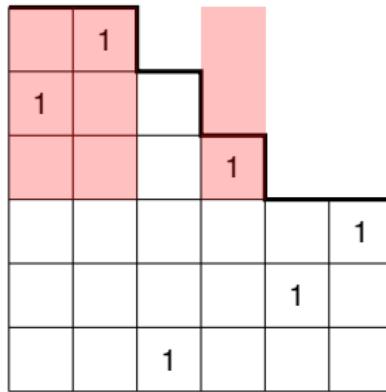
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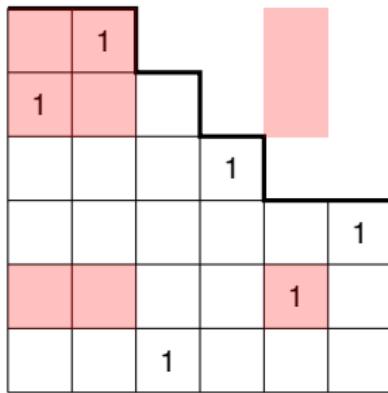
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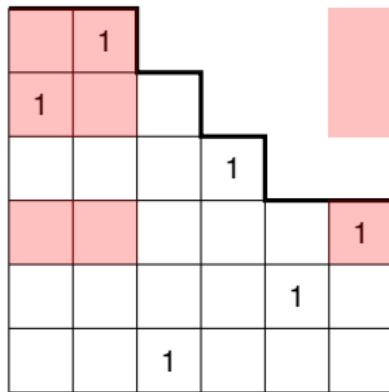
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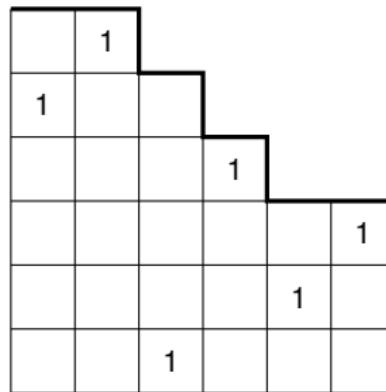
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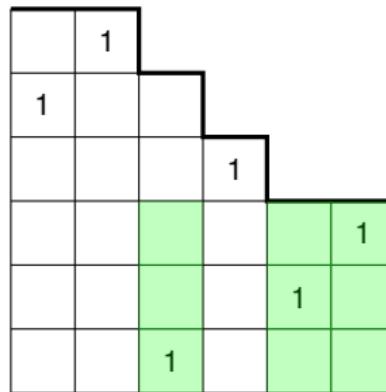


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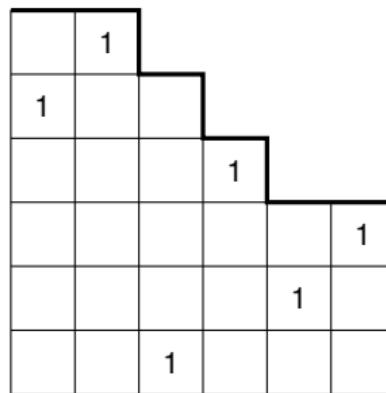
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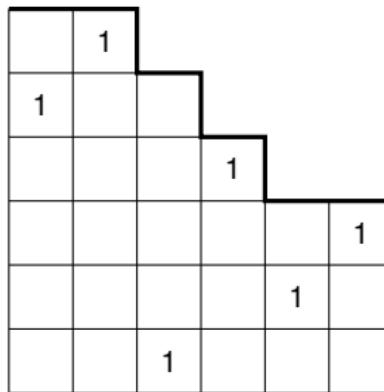
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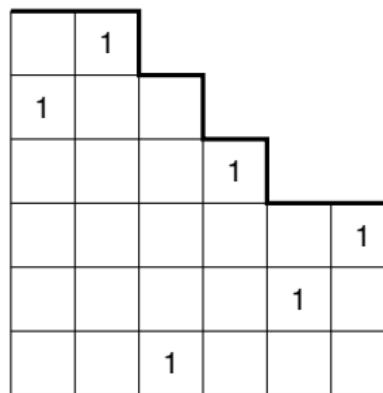
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- Patterns σ and τ are **NE-shape-Wilf-equivalent**, denoted $\sigma \sim_s \tau$, if for any fixed NE-shape Λ , equal number of **traversals** (or **transversals**) of Λ avoid σ and τ .

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- SW-shape-Wilf-equivalence can be defined similarly, denoted $\sigma \curvearrowleft_s \tau$

Previous results for NE-shape-Wilf equivalence

- Direct sum \oplus and skew-sum \ominus :

$$\sigma \oplus \tau = \begin{array}{c} \boxed{\tau} \\ \hline \boxed{\sigma} \end{array} \qquad \sigma \ominus \tau = \begin{array}{c} \boxed{\sigma} \\ \hline \boxed{\tau} \end{array}$$

- Backelin, West, Xin, 2007:

- $S' \sim_s S'' \implies S' \oplus S \sim_s S'' \oplus S$, for any sets of patterns S, S', S''
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- $I_n \sim_s J_n$ and $I_n \sim_s J_n$, for the identity I_n and anti-identity J_n patterns of any size $n \geq 0$

(Theorem 3.1, Corollary 3.2, Corollary 3.3, Corollary 3.4, Corollary 3.5, Corollary 3.6)

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- $(1, 3)_3 \sim_s (2, 3)_3 \sim_s (3, 3)_3 \sim_s (3, 1)_3 \sim_s (2, 1)_3 \sim_s (1, 2)_3$
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- $S' \sim_s S'' \implies S' \oplus S \sim_s S'' \oplus S$, for any sets of patterns S, S', S''
- Equivalently, $S' \sim_s S'' \implies S \oplus S' \sim_s S \oplus S''$, for any sets of patterns S, S', S''
- $I_n \sim_s J_n$ and $I_n \sim_s J_n$, for the identity I_n and anti-identity J_n patterns of any size $n \geq 0$
 - This follows from iterating $J_{n+1} = J_n \ominus 1 = 1 \ominus J_n \sim_s J_n \oplus 1$ for any $n \geq 0$

- Bloom, Elizalde, 2014:

- $(1, 3)_3 \sim_s (2, 3)_3 \sim_s (3, 3)_3 \sim_s (3, 1)_3 \sim_s (2, 1)_3 \sim_s (1, 2)_3$
- Equivalently, $(3, 1)_3 \sim_s (2, 1)_3 \sim_s (1, 1)_3 \sim_s (1, 3)_3 \sim_s (3, 2)_3 \sim_s (2, 3)_3$

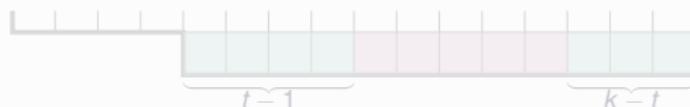
Results I

Theorem (B.-Han-Kitaev-Zhang, 2024)

$(t, 1)_k \rightsquigarrow_s (1, 1)_k$ for all $1 \leq t \leq k$, so all $(t, 1)$, $t \geq 1$, are SW-shape-Wilf-equivalent.

Proof.

Grow the shape by adding the new bottom row and the column of the new bottom 1.



The number of possible insertion cells in a bottom row of length ℓ is $\min(\ell, k - 1)$. □

This result was originally stated in terms of partially ordered patterns (POPs).

Corollary

$(1t, 12)_k \rightsquigarrow_s (12, 12)_k$ for all $2 \leq t \leq k$, so all $(1t, 12)$, $t \geq 2$, are SW-shape-Wilf-equivalent.

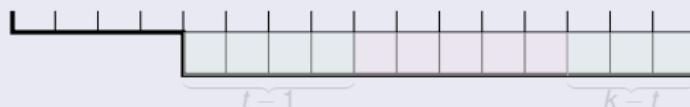
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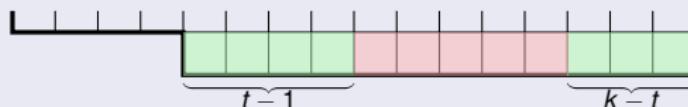
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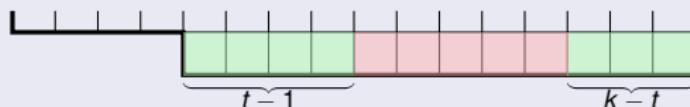
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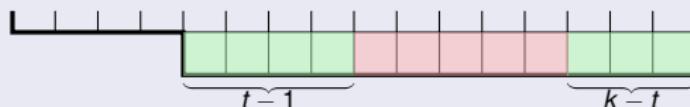
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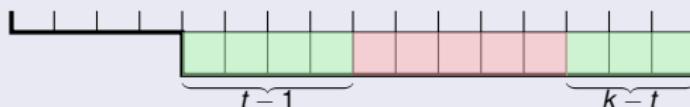
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Results II

Theorem

$((t, t+1), (1, 2))_k \sim (12, 12)_k$ for all $1 \leq t \leq k-1$, so all $((t, t+1), (1, 2)), t \geq 1$, are Wilf-equivalent.

Moreover, for $k \geq 3$, $|\text{Av}_n(((t, t+1), (1, 2))_k)| = \begin{cases} n!, & \text{if } n < k-3, \\ (k-3)!r_{k-3}(n), & \text{if } n \geq k-3, \end{cases}$,

where $r_{k-3}(n)$ is the n -th $(k-3)$ -Schröder number, the number of Schröder paths from $(0, 0)$ to $(2n, 0)$ on or above the x -axis with steps $U = (1, 1)$, $D = (1, -1)$, and steps $H = (2, 0)$ of $k-3$ colors.

Proof Sketch.

For $\sigma \in \text{Av}_n(((t, t+1), (1, 2))_k)$, consider the top $k-2$ values of σ (their order is irrelevant) and the blocks into which they split σ . Count the number of the points in the region that avoids 12. Use this to find a functional equation for the generating function with 1 auxiliary variable, then use the kernel method. □

Proof Sketch (cont'd)

Let $i = t - 1, j = k - t - 1$ (so $k - 2 = i + j$), and let $R_{ij} = ((i+1, i+2), (1, 2))_{i+j+2}$. Note that $R_{00} = (12, 12)_2 = 12$.



$$F(x, y) = 1 + (k-2)x \frac{yF(x, y) - F(x, 1)}{y-1} + xyF(x, y)$$

$$(y-1-(k-3)xy-x)^2 F(x, y) = y - (1 + (k-2)xF(x, 1))$$

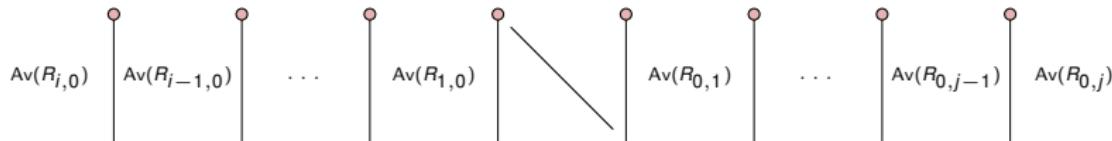
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OGF over $j \geq k-3$:

$$(k-3)!x^{k-3} + (k-2)!x^{k-2}F(x, 1) = (k-3)!x^{k-3}(1 + (k-2)xF(x, 1)) = (k-3)!x^{k-3}R_{k-3}(x).$$

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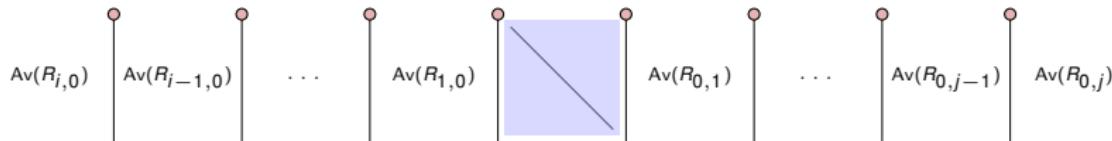
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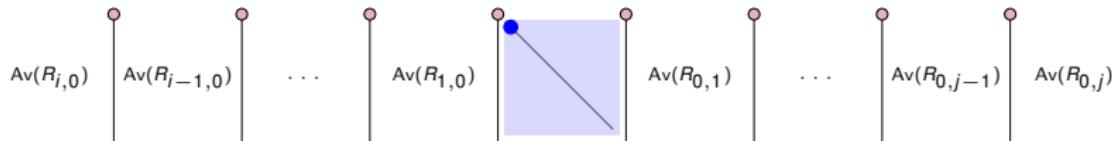
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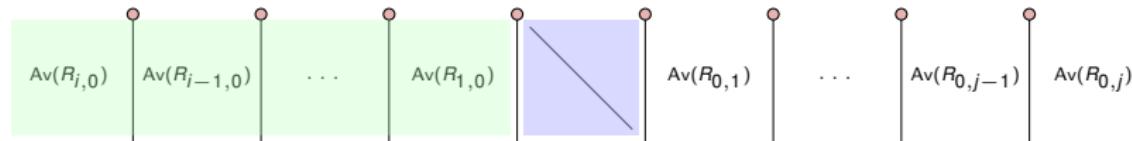
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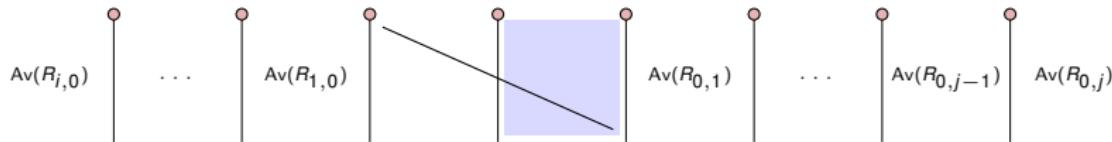
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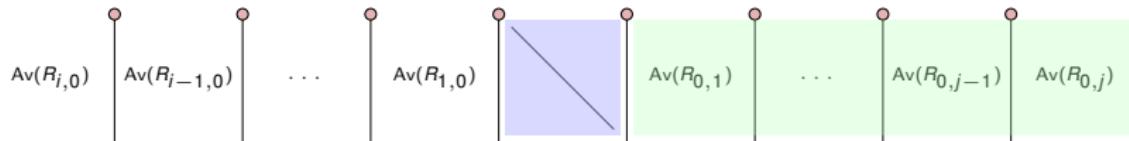
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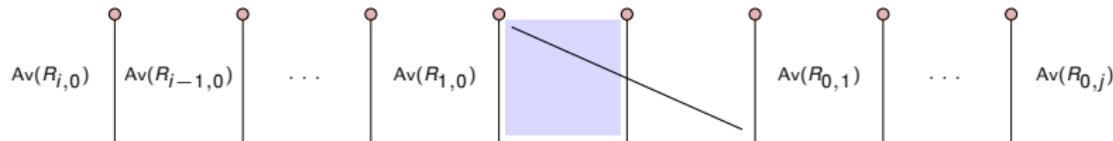
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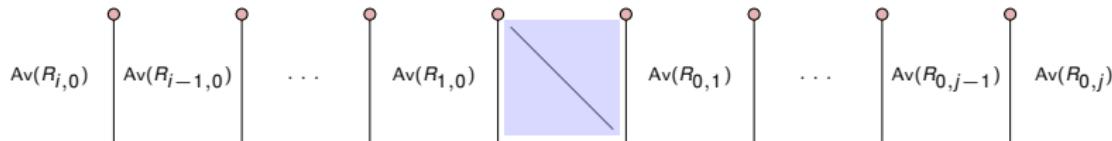
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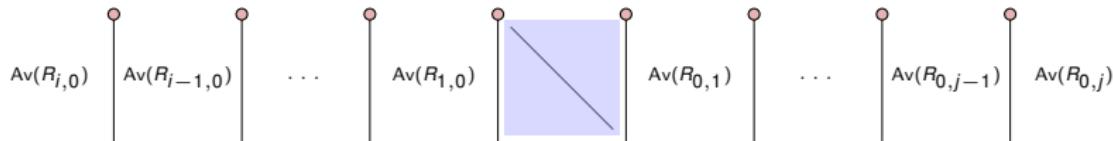
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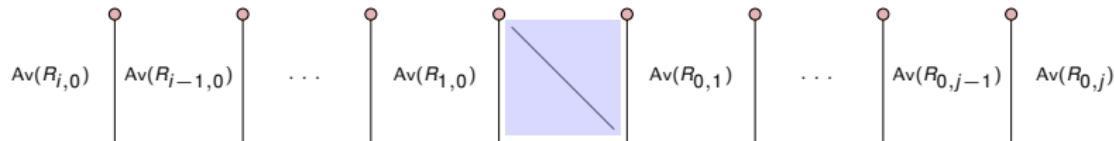
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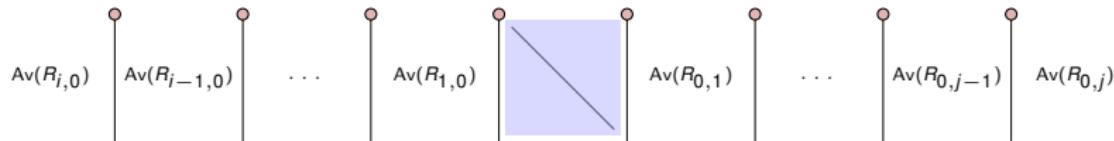
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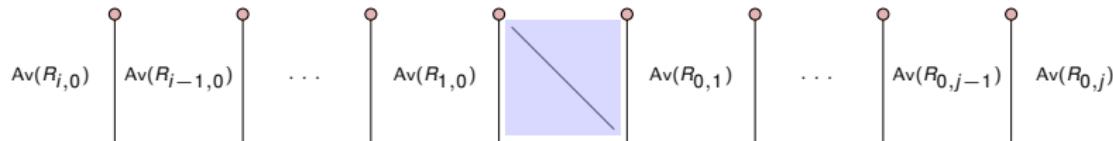
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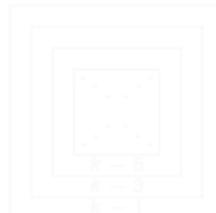
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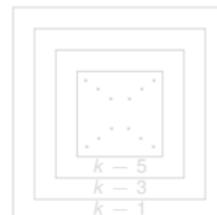
Conjectures

- $(ij, 12)$ are Wilf-equivalent for all (i, j) such that $i = 1$ or $j = i \pm 1$ or $j = i \pm 2$.
- (BHKZ'24): $(k, 2)_k \rightsquigarrow_s (k, 1)_k$, and thus $(1k, 13)_k \rightsquigarrow_s (1k, 12)_k$ for all $k \geq 3$.
 - Proved by Wang and Yan (2025).
- For each $i \geq 1$, $((j+1, j+2, \dots, j+i), (1, 2, \dots, i))$ are Wilf-equivalent for all $j \geq 0$.
- $(135, \sigma)$ are 5-Wilf-equivalent for all $\sigma \in S_3$.
Equivalently, $(14253, 15243) \sim (14352, 15342) \sim (24153, 25143)$.
- $(24, 12)_4 \rightsquigarrow_s (24, 21)_4$. Equivalently, $(3142, 4132) \rightsquigarrow_s (3241, 4231)$.
- $(134, 123) \sim (234, 213)$. Equivalently, $(123, 134) \sim (123, 324)$.
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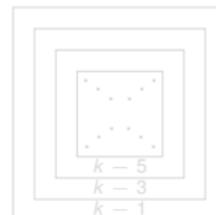
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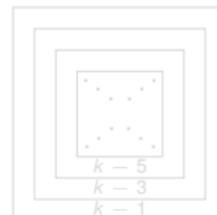
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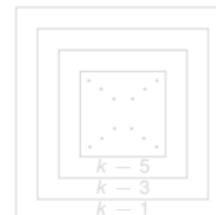
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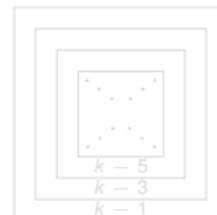
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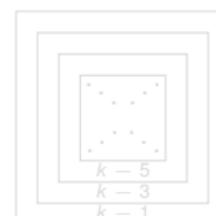
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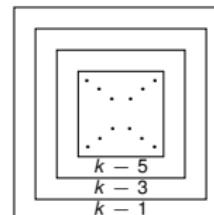
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Examples

We know that $|Av_n((t, 1)_k)| = (k - 1)!(\textcolor{red}{k - 1})^{n-k+1}$ for $1 \leq t \leq k$.

Here are more examples for (i, j) with $2 \leq i, j \leq k - 1$.

- $|Av_n((2, 2)_4)| = A128445(n) = 4((n-2)^2 + 1)$ for $n \geq 5$.
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