Descents of permutations with only even or only odd cycles

Sergi Elizalde

Dartmouth

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- in one-line notation, e.g. $\pi=3\,1\,7\,6\,5\,4\,2$,
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Example

$$\begin{array}{l} \mathcal{S}_3^o = \{(1,2,3), (1,3,2), (1)(2)(3)\}, \\ \mathcal{S}_3^e = \{(1,2)(3), (1,3)(2), (2,3)(1)\}. \end{array}$$

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• Write each cycle of $\pi \in \mathcal{S}_n^o$ starting with its largest element, and order the cycles by increasing first element, e.g. $\pi = (4)(5,1,3)(7,2,6)(8)$.

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- Write each cycle of $\pi \in \mathcal{S}_n^o$ starting with its largest element, and order the cycles by increasing first element, e.g. $\pi = (4)(5,1,3)(7,2,6)(8)$.
- Move the last element of the 1st cycle to the end of the 2nd cycle, the last element of the 3rd cycle to the end of the 4th cycle, etc., e.g.

$$(\mathbf{4})(5,1,3)(7,2,\mathbf{6})(8) \mapsto (5,1,3,\mathbf{4})(7,2)(8,\mathbf{6}) \in \mathcal{S}_n^e.$$

A surprising refinement

Definition

For $\pi = \pi_1 \pi_2 \dots \pi_n \in \mathcal{S}_n$,

- its descent set is $Des(\pi) = \{i : \pi_i > \pi_{i+1}\},\$
- its ascent set is $Asc(\pi) = \{i : \pi_i < \pi_{i+1}\}.$

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Theorem (Adin, Hegedűs, Roichman '25)

For any n and any subset $J \subseteq [n-1]$,

$$|\{\pi \in \mathcal{S}^o_n : \mathrm{Asc}(\pi) = J\}| = |\{\pi \in \mathcal{S}^e_n : \mathrm{Des}(\pi) = J\}|.$$

A surprising refinement

$$|\{\pi \in \mathcal{S}_n^o : \operatorname{Asc}(\pi) = J\}| = |\{\pi' \in \mathcal{S}_n^e : \operatorname{Des}(\pi') = J\}|$$

Example (n=4)

$\pi \in \mathcal{S}_4^{o}$	$Asc(\pi) = Des(\pi')$	$\pi'\in\mathcal{S}_4^{e}$
(1,2,4)(3) = 2431	{1}	(1,4,3,2) = 4123
(1,4,2)(3) = 4132	{2}	(1,2,4,3) = 2413
(1,3,4)(2) = 3241	{2}	(1,3)(2,4) = 3412
(1,4,3)(2) = 4213	{3}	(1,2,3,4) = 2341
(2,3,4)(1) = 1342	$\{1, 2\}$	(1,4,2,3) = 4312
(2,4,3)(1) = 1423	$\{1, 3\}$	(1,3,4,2) = 3142
(1,2,3)(4) = 2314	$\{1, 3\}$	(1,2)(3,4) = 2143
(1,3,2)(4) = 3124	$\{2, 3\}$	(1,3,2,4) = 3421
(1)(2)(3)(4) = 1234	$\{1, 2, 3\}$	(1,4)(2,3) = 4321

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Unfortunately, Bóna's bijection $\mathcal{S}_n^o \to \mathcal{S}_n^e$ does not behave well with respect to Asc and Des .

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Our goal is to provide a bijective proof.

Unfortunately, Bóna's bijection $\mathcal{S}_n^o \to \mathcal{S}_n^e$ does not behave well with respect to Asc and Des .

For any n and any subset $S \subseteq [n-1]$, we will construct an explicit bijection

$$\{\pi \in \mathcal{S}_n^o : \operatorname{Asc}(\pi) \subseteq S\} \longleftrightarrow \{\pi \in \mathcal{S}_n^e : \operatorname{Des}(\pi) \subseteq S\}.$$

Structure of the bijection

$$\{\pi \in \mathcal{S}_n^o : \operatorname{Asc}(\pi) \subseteq S\}$$



$$\{\pi \in \mathcal{S}_n^e : \mathrm{Des}(\pi) \subseteq \mathcal{S}\}$$

Structure of the bijection

$$\{\pi \in \mathcal{S}_n^o : \mathrm{Asc}(\pi) \subseteq \mathcal{S}\}$$

Multisets of odd, distinct necklaces

$$\{\pi \in \mathcal{S}_n^e : \mathrm{Des}(\pi) \subseteq S\}$$

$$\downarrow$$

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Multisets of odd, distinct necklaces

Words whose Lyndon factors are odd and distinct



$$\{\pi \in \mathcal{S}_n^e : \mathrm{Des}(\pi) \subseteq \mathcal{S}\}$$

Multisets of even necklaces

|||

Words whose Lyndon factors are even

Let $A = \{a, b, c, \dots\}$, $\mathcal{W} = \text{finite words over } A$.

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A word $u \in \mathcal{W}$ is primitive if it is not of the form $u = r^j$ for $j \ge 2$.

Example: ababa is primitive, but $abab = (ab)^2$ is not.

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For a set
$$S = \{s_1, s_2, \dots, s_k\} \subseteq [n-1]$$
 with $s_1 < \dots < s_k$, let $\alpha = (s_1, s_2 - s_1, \dots, n - s_k)$ and $\operatorname{wt}(S) = \operatorname{a}^{\alpha_1} \operatorname{b}^{\alpha_2} \operatorname{c}^{\alpha_3} \dots$

Example: If $S = \{2, 3\} \subseteq [5]$, then $\alpha = (2, 1, 3)$ and $wt(S) = a^2bc^3$.

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$$\Phi_{\mathcal{S}}: \{\pi \in \mathcal{S}_n : \mathrm{Des}(\pi) \subseteq \mathcal{S}\} \to \{M \in \mathcal{M}_n : \mathrm{wt}(M) = \mathrm{wt}(\mathcal{S})\}$$

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Let n = 8 and $S = \{4, 7\}$, so $wt(S) = a^4b^3c$.

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To recover π from the multiset of necklaces:

 Replace each entry with the periodic sequence obtained by reading around the necklace.

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To recover π from the multiset of necklaces:

- Replace each entry with the periodic sequence obtained by reading around the necklace.
- Order these sequences lexicographically (breaking ties consistently).

Example

The multiset of necklaces (a,b)(a,b)(a,a,b,c) gives periodic sequences (abab...,baba...)(abab...,baba...)(aabc...,abca...,bcaa...,caab...). We get $\pi = (3,6)(2,5)(1,4,7,8)$.

 $\mathcal{M}_n^e = \text{multisets of necklaces of even length,}$ except possibly for one necklace of length one.

$$\{\pi \in \mathcal{S}_n^o : \operatorname{Asc}(\pi) \subseteq S\}$$



$$\{\pi \in \mathcal{S}_n^e : \mathrm{Des}(\pi) \subseteq S\}$$

$$\downarrow \Phi_S$$

$$\{M \in \mathcal{M}_n^e : \mathrm{wt}(M) = \mathrm{wt}(S)\}$$

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 $\mathcal{M}_{p}^{o} = \text{multisets of distinct necklaces of odd length.}$

We get a bijection

$$\Xi_{S}: \{\pi \in \mathcal{S}_{n}^{o}: \mathrm{Asc}(\pi) \subseteq S\} \rightarrow \{M \in \mathcal{M}_{n}^{o}: \mathrm{wt}(M) = \mathrm{wt}(S)\}.$$

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We will interpret multisets of necklaces as words.

A Lyndon word is a primitive word that is lexicographically smaller than all of its cyclic rotations. Denote the set of Lyndon words by $\mathcal{L} \subseteq \mathcal{W}$.

Example: $aabab \in \mathcal{L}$, but $ababa \notin \mathcal{L}$.

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Theorem (Lyndon '58)

Every $w \in \mathcal{W}$ has a unique Lyndon factorization $w = \ell_1 | \ell_2 | \dots | \ell_m$ where $\ell_i \in \mathcal{L}$ for all i, and $\ell_1 \geq \ell_2 \geq \dots \geq \ell_m$ lexicographically.

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w = dedccedcdbdbdaabd

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Example

 $\begin{aligned} \textit{w} &= \texttt{dedccedcdbdbdaabd} = \texttt{de}|\texttt{d}|\texttt{ccedcd}|\texttt{bd}|\texttt{bd}|\texttt{aabd} \\ &\leftrightarrow (\texttt{d},\texttt{e})(\texttt{d})(\texttt{c},\texttt{c},\texttt{e},\texttt{d},\texttt{c},\texttt{d})(\texttt{b},\texttt{d})(\texttt{b},\texttt{d})(\texttt{a},\texttt{a},\texttt{b},\texttt{d}). \end{aligned}$

We identify multisets of necklaces with words.

Define the following sets of length-*n* words:

 \mathcal{W}_n^e = words all of whose Lyndon factors have even length, except possibly for one factor which has length one.

 $\mathcal{W}_n^o = \text{words all of whose Lyndon factors have odd length and are distinct.}$

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$$\{\pi \in \mathcal{S}_{n}^{o} : \operatorname{Asc}(\pi) \subseteq S\} \qquad \qquad \{\pi \in \mathcal{S}_{n}^{e} : \operatorname{Des}(\pi) \subseteq S\}$$

$$\Xi_{S} \downarrow \qquad \qquad \downarrow \Phi_{S}$$

$$\{M \in \mathcal{M}_{n}^{o} : \operatorname{wt}(M) = \operatorname{wt}(S)\} \qquad \qquad \{M \in \mathcal{M}_{n}^{e} : \operatorname{wt}(M) = \operatorname{wt}(S)\}$$

$$\parallel \qquad \qquad \parallel \qquad \qquad \parallel \parallel$$

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$$\| \{W \in \mathcal{W}_{n}^{o} : \operatorname{wt}(w) = \operatorname{wt}(S)\}$$

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$$\{W \in \mathcal{W}_{n}^{e} : \operatorname{wt}(w) = \operatorname{wt}(S)\}$$

We want a weight-preserving bijection between \mathcal{W}_n^o and \mathcal{W}_n^e .

Given $w \in \mathcal{W}_n^o$ (suppose *n* is even), initially set O = w and E = -.

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Repeat until O is empty:

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Repeat until O is empty:

- Let $O = o_1 | o_2 | \dots | o_m$ be the Lyndon factorization of O.
- If $|o_m| \ge 2$, write $o_m = r_1 s$ where s is its lexicographically smallest proper suffix.

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- Update (*O*, *E*) to

$$\begin{cases} (o_1o_2\dots o_{m-1}r,\,sE) & \text{if } o_m \text{ is splittable and } |s| \text{ is even},\\ (o_1o_2\dots o_{m-1}s,\,rE) & \text{if } o_m \text{ is splittable and } |r| \text{ is even},\\ (o_1o_2\dots o_{m-2},\,o_mo_{m-1}E) & \text{if } o_m \text{ is not splittable}. \end{cases}$$

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Repeat until O is empty:

- Let $O = o_1 | o_2 | \dots | o_m$ be the Lyndon factorization of O.
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Set $\Psi(w) = E$.

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Example

$$w = dadcdebccc$$

- Let $O = o_1 | o_2 | \dots | o_m$ be the Lyndon factorization of O.
- If $|o_m| \ge 2$, write $o_m = r$'s where s is its smallest proper suffix.
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Example

$$O$$
 E $w = d | adcdebccc -$

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Example

$$O$$
 E $w = d \frac{\text{adcde!bccc}}{}$

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Example

$$egin{array}{ccc} O & E \ \hline w = & d | adcde'_bccc & - \ & d | adcde & bccc \ \hline \end{array}$$

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Example

$$O$$
 $W = d \mid adcdebccc$
 $d \mid adcde$
 $bccc$

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 $W = d | adcdebccc$
 $d | ad | cde$
 $d | bccc$

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Example

	0	Е
w =	d adcdebccc	_
	d <mark>ad</mark> ¦cde	bccc
	d cde	<u>ad</u> bccc

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w =	d adcdebccc	_	
	d adcde	bccc	
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	0	Ε	
w =	d adcdebccc	_	
	d adcde	bccc	
	d c¦de	adbccc	
	_	cdedadbccc	

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w =	d adcdebccc	_	
	d adcde	bccc	
	d cde	adbccc	
	_	cdedadbccc	$=\Psi(w)$

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Example

•

	0	Ε	
w =	d adcdebccc	_	
	d adcde	bccc	
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	_	cded adbccc	$=\Psi(w)$

Theorem

The map $\Psi: \mathcal{W}_n^o \to \mathcal{W}_n^e$ is a weight-preserving bijection.

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The map $\Psi: \mathcal{W}_n^o \to \mathcal{W}_n^e$ is a weight-preserving bijection.

Composing the three maps, we obtain the desired bijection:

$$\{\pi \in \mathcal{S}_{n}^{o} : \operatorname{Asc}(\pi) \subseteq S\}$$

$$\Xi_{S} \downarrow \qquad \qquad \{\pi \in \mathcal{S}_{n}^{e} : \operatorname{Des}(\pi) \subseteq S\}$$

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$$\{W \in \mathcal{W}_{n}^{o} : \operatorname{wt}(W) = \operatorname{wt}(S)\}$$

$$\{W \in \mathcal{W}_{n}^{e} : \operatorname{wt}(W) = \operatorname{wt}(S)\}$$

$$|\{w\in\mathcal{W}^o_n:\operatorname{wt}(w)=x_1^{\alpha_1}x_2^{\alpha_2}\dots\}|=|\{w\in\mathcal{W}^e_n:\operatorname{wt}(w)=x_1^{\alpha_1}x_2^{\alpha_2}\dots\}|$$

$$|\{w \in \mathcal{W}_n^o : \operatorname{wt}(w) = x_1^{\alpha_1} x_2^{\alpha_2} \dots\}| = |\{w \in \mathcal{W}_n^e : \operatorname{wt}(w) = x_1^{\alpha_1} x_2^{\alpha_2} \dots\}|$$

Proof:

ullet Write $\mathcal{L}=\mathcal{L}^o\sqcup\mathcal{L}^e$, separating Lyndon words of odd and even length.

$$|\{w \in \mathcal{W}_n^o : \operatorname{wt}(w) = x_1^{\alpha_1} x_2^{\alpha_2} \dots\}| = |\{w \in \mathcal{W}_n^e : \operatorname{wt}(w) = x_1^{\alpha_1} x_2^{\alpha_2} \dots\}|$$

Proof:

- ullet Write $\mathcal{L}=\mathcal{L}^o\sqcup\mathcal{L}^e$, separating Lyndon words of odd and even length.
- GF for $\mathcal{W}^o = \bigcup_{n \geq 0} \mathcal{W}^o_n$ (words with odd and distinct Lyndon factors):

$$\sum_{\mathbf{w} \in \mathcal{W}^{o}} \operatorname{wt}(\mathbf{w}) = \prod_{\ell \in \mathcal{L}^{o}} (1 + \operatorname{wt}(\ell)).$$

$$|\{w \in \mathcal{W}_n^o : \operatorname{wt}(w) = x_1^{\alpha_1} x_2^{\alpha_2} \dots\}| = |\{w \in \mathcal{W}_n^e : \operatorname{wt}(w) = x_1^{\alpha_1} x_2^{\alpha_2} \dots\}|$$

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• GF for $\mathcal{W}^e = \bigcup_{n \geq 0} \mathcal{W}^e_n$ (even Lyndon factors, except one of length 1):

$$\sum_{\boldsymbol{w} \in \mathcal{W}^e} \operatorname{wt}(\boldsymbol{w}) = (1 + x_1 + x_2 + \dots) \prod_{\ell \in \mathcal{L}^e} \frac{1}{1 - \operatorname{wt}(\ell)}.$$

$$|\{w \in \mathcal{W}_n^o : \operatorname{wt}(w) = x_1^{\alpha_1} x_2^{\alpha_2} \dots\}| = |\{w \in \mathcal{W}_n^e : \operatorname{wt}(w) = x_1^{\alpha_1} x_2^{\alpha_2} \dots\}|$$

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$$\sum_{\mathbf{w}\in\mathcal{W}^e}\operatorname{wt}(\mathbf{w})=(1+x_1+x_2+\dots)\prod_{\ell\in\mathcal{L}^e}\frac{1}{1-\operatorname{wt}(\ell)}.$$

Thus, what we want to prove is

$$\prod_{\ell \in \mathcal{L}^o} \left(1 + \operatorname{wt}(\ell)\right) = \left(1 + x_1 + x_2 + \dots\right) \prod_{\ell \in \mathcal{L}^e} \frac{1}{1 - \operatorname{wt}(\ell)},$$

or equivalently,

$$\prod_{\ell \in \mathcal{L}^o} (1 + \operatorname{wt}(\ell)) \prod_{\ell \in \mathcal{L}^e} (1 - \operatorname{wt}(\ell)) = 1 + x_1 + x_2 + \dots$$

or equivalently,

$$\prod_{\ell \in \mathcal{L}^o} (1 + \operatorname{wt}(\ell)) \prod_{\ell \in \mathcal{L}^e} (1 - \operatorname{wt}(\ell)) = 1 + x_1 + x_2 + \dots$$

• Substituting $x_i \mapsto -x_i$ for all i, this is equivalent to

$$\prod_{\ell \in \mathcal{L}^o} (1 - \operatorname{wt}(\ell)) \prod_{\ell \in \mathcal{L}^e} (1 - \operatorname{wt}(\ell)) = 1 - x_1 - x_2 - \dots,$$

or equivalently,

$$\prod_{\ell \in \mathcal{L}^o} (1 + \operatorname{wt}(\ell)) \prod_{\ell \in \mathcal{L}^e} (1 - \operatorname{wt}(\ell)) = 1 + x_1 + x_2 + \dots$$

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or

$$\prod_{\ell \in \mathcal{L}} \frac{1}{1 - \operatorname{wt}(\ell)} = \frac{1}{1 - x_1 - x_2 - \dots}.$$

or equivalently,

$$\prod_{\ell \in \mathcal{L}^o} (1 + \operatorname{wt}(\ell)) \prod_{\ell \in \mathcal{L}^e} (1 - \operatorname{wt}(\ell)) = 1 + x_1 + x_2 + \dots$$

• Substituting $x_i \mapsto -x_i$ for all i, this is equivalent to

$$\prod_{\ell \in \mathcal{L}^o} (1 - \operatorname{wt}(\ell)) \prod_{\ell \in \mathcal{L}^e} (1 - \operatorname{wt}(\ell)) = 1 - x_1 - x_2 - \dots,$$

or

$$\prod_{\ell \in \mathcal{L}} \frac{1}{1 - \operatorname{wt}(\ell)} = \frac{1}{1 - x_1 - x_2 - \dots}.$$

• But this holds because every word has a unique Lyndon factorization!

